IN THE CLAIMS

Please replace all prior claims in the present application with the following claims, in which claims 1, 18, and 35 are currently amended.

1. (Currently Amended) A method of performing bandwidth allocations, the method comprising:

receiving a bandwidth request from a terminal;

determining bandwidth request type and priority of the received bandwidth request;

placing the bandwidth request in one of a plurality of a global queues based upon the determining step, each of the global queues corresponding to a data rate of each of a plurality of channels;

moving the bandwidth request from the one global queue to one of a plurality of local queues, the plurality of local queues corresponding to the plurality of channels, wherein the bandwidth request is moved based on loading of the channels; and

allocating the transmission slots in response to the bandwidth request stored in the one local queue.

2. (Original) The method as in claim 1, wherein the bandwidth request in the receiving step is at least one of a rate request and a volume request, the rate request specifying a constant number of transmission slots, the volume request specifying a specific number of transmission slots.

3. (Original) The method as in claim 1, wherein the bandwidth request in the receiving step is a rate request, the method further comprising:

filling the one local queue with subsequent rate requests up to a queuing threshold; and

filling another one of the local queues with additional rate requests upon filling the one local queue beyond the queuing threshold.

- 4. (Original) The method as in claim 3, wherein the queuing threshold in the step of filling the one local queue is predetermined.
- 5. (Original) The method as in claim 3, wherein the queuing threshold in the step of filling the one local queue is dynamically established according to a total number of rate requests in the plurality of local queues.
- 6. (Original) The method as in claim 1, wherein the global queues in the placing step are designated according to the bandwidth request type and the associated priority.
- 7. (Original) The method as in claim 1, wherein the bandwidth request type and priority of the received bandwidth request include a high priority rate request, a low priority rate request, a high priority volume request, and a low priority volume request.

8. (Original) The method as in claim 1, wherein the bandwidth request in the receiving step is a volume request and is received over at least one of a contention channel and a data channel, the method further comprising:

receiving a piggybacked volume request over the data channel;

placing the piggybacked volume request in a corresponding one of the global queues;

determining whether the plurality of channels are oversubscribed; and selectively discarding the piggybacked volume request based upon the step of determining whether the plurality of channels are oversubscribed.

9. (Original) The method as in claim 8, wherein the step of determining whether the plurality of channels are oversubscribed comprises:

determining whether each of the plurality of local queues exceeds a respective queuing threshold.

10. (Original) The method as in claim 1, wherein the plurality of channels are designated as data channels that are sequentially ordered, the allocating step comprising:

selectively assigning the transmission slots according to a prescribed order of the data channels based upon the bandwidth request type.



- 11. (Original) The method as in claim 10, wherein the prescribed order in the selectively assigning step begins with the first data channel if the bandwidth request type is rate request.
- 12. (Original) The method as in claim 10, wherein the prescribed order in the selectively assigning step begins with the last data channel if the bandwidth request type is volume request.
 - 13. (Original) The method as in claim 1, further comprising: receiving a plurality of rate requests;
 receiving a defragmentation command; and

moving the rate requests from the local queues to the corresponding global queues for reallocation in response to the defragmentation command.

14. (Original) The method as in claim 1, wherein the bandwidth request is an original volume request, the method further comprising:

receiving a follow-up volume request;

associating the follow-up volume request with the original volume request; and placing the follow-up volume request to a particular local queue that stored the original bandwidth request among the plurality of local queues.

15. (Original) The method as in claim 14, wherein the associating step comprises:



maintaining a database of pointers for the terminal, one of the pointers specifying the particular local queue.

16. (Original) The method as in claim 1, further comprising: receiving a plurality of volume requests; and spreading the volume requests across each of the local queues.

17. (Original) The method as in claim 16, wherein each of the local queues has a counter that counts a quantity of the volume requests in the respective local queue, the distributing step comprising:

comparing counter values of the counters with respective predetermined thresholds corresponding to the local queues.

18. (Currently Amended) A communication system for performing bandwidth allocations, the system comprising:

a plurality of global queues, each of the global queues being configured to store a bandwidth request received from a terminal;

a bandwidth control processor (BCP) coupled the plurality of global queues, the bandwidth control processor being configured to determine bandwidth request type and priority of the received bandwidth request and to place the bandwidth request in one of a plurality of a global queues based upon the determined bandwidth request type and priority, wherein each of the global queues corresponds to a data rate of each of a plurality of channels; and

a plurality of local queues coupled to the BCP, the plurality of local queues corresponding to the plurality of channels, one of the plurality of local queues storing the bandwidth request <u>is</u> moved from the one global queue <u>based on loading of the channels</u>,

wherein the BCP allocates the transmission slots in response to the bandwidth request stored in the one local queue.

- 19. (Original) The system as in claim 18, wherein the bandwidth request is at least one of a rate request and a volume request, the rate request specifying a constant number of transmission slots, the volume request specifying a specific number of transmission slots.
- 20. (Original) The system as in claim 18, wherein the bandwidth request is a rate request, the one local queue being filled with subsequent rate requests up to a queuing threshold, another one of the local queues being filled up with additional rate requests in response to the one local queue being filled beyond the queuing threshold.
- 21. (Original) The system as in claim 20, wherein the queuing threshold is predetermined.
- 22. (Original) The system as in claim 20, wherein the queuing threshold is dynamically established according to a total number of rate requests in the plurality of local queues.



- 23. (Original) The system as in claim 18, wherein the global queues are designated according to the bandwidth request type and the associated priority.
- 24. (Original) The system as in claim 18, wherein the bandwidth request type and priority of the received bandwidth request include a high priority rate request, a low priority rate request, a high priority volume request, and a low priority volume request.
- 25. (Original) The system as in claim 18, wherein the bandwidth request is a volume request and is received over at least one of a contention channel and a data channel, a follow-up volume request associated with the volume request being placed in a corresponding one of the global queues, the BCP selectively discarding the follow-up volume request upon determining that the plurality of channels are oversubscribed.
- 26. (Original) The system as in claim 25, wherein the BCP determines oversubscription of the plurality of channels by examining whether each of the plurality of local queues exceeds a respective queuing threshold.
- 27. (Original) The system as in claim 18, wherein the plurality of channels are designated as data channels that are sequentially ordered, the BCP selectively assigning the transmission slots according to a prescribed order of the data channels based upon the bandwidth request type.

Month of the second of the sec

- 28. (Original) The system as in claim 27, wherein the prescribed order begins with the first data channel if the bandwidth request type is rate request.
- 29. (Original) The system as in claim 27, wherein the prescribed order begins with the last data channel if the bandwidth request type is volume request.
- 30. (Original) The system as in claim 18, wherein the BCP is configured to move the rate requests from the local queues to the corresponding global queues for reallocation in response to a defragmentation command.
- 31. (Original) The system as in claim 18, wherein the bandwidth request is an original volume request, the BCP associating a follow-up volume request with the original volume request and placing the follow-up volume request to a particular local queue that stored the original bandwidth request among the plurality of local queues.
- 32. (Original) The system as in claim 31, further comprising:

 a database coupled to the BCP, the database storing a plurality of pointers for the terminal, one of the pointers specifying the particular local queue.
- 33. (Original) The system as in claim 18, wherein one of the global queues stores a plurality of volume requests, the BCP spreading the volume requests across the local queues.

34. (Original) The system as in claim 33, wherein each of the local queues has a counter that counts a quantity of the volume requests in the respective local queue, the BCP comparing counter values of the counters with predetermined thresholds corresponding to the local queues.

35. (Currently Amended) A computer readable medium containing program instructions for execution on a computer system, which when executed by a computer, cause the computer system to perform method steps for allocating bandwidth, said method comprising the steps of:

receiving a bandwidth request from a terminal;

determining bandwidth request type and priority of the received bandwidth request;

placing the bandwidth request in one of a plurality of a global queues based upon the determining step, each of the global queues corresponding to a data rate of each of a plurality of channels;

moving the bandwidth request from the one global queue to one of a plurality of local queues, the plurality of local queues corresponding to the plurality of channels, wherein the bandwidth request is moved based on loading of the channels; and

allocating the transmission slots in response to the bandwidth request stored in the one local queue.



- 36. (Original) The computer-readable medium as in claim 35, wherein the bandwidth request in the receiving step is at least one of a rate request and a volume request, the rate request specifying a constant number of transmission slots, the volume request specifying a specific number of transmission slots.
- 37. (Original) The computer-readable medium as in claim 35, wherein the bandwidth request in the receiving step is a rate request, the computer-readable medium further comprising computer-executable instructions for causing the computer system to perform the steps of:

filling the one local queue with subsequent rate requests up to a queuing threshold; and

filling another one of the local queues with additional rate requests upon filling the one local queue beyond the queuing threshold.

- 38. (Original) The computer-readable medium as in claim 37, wherein the queuing threshold in the step of filling the one local queue is predetermined.
- 39. (Original) The computer-readable medium as in claim 37, wherein the queuing threshold in the step of filling the one local queue is dynamically established according to a total number of rate requests in the plurality of local queues.

Moy.

- 40. (Original) The computer-readable medium as in claim 35, wherein the global queues in the placing step are designated according to the bandwidth request type and the associated priority.
- 41. (Original) The computer-readable medium as in claim 35, wherein the bandwidth request type and priority of the received bandwidth request include a high priority rate request, a low priority rate request, a high priority volume request, and a low priority volume request.
- 42. (Original) The computer-readable medium as in claim 35, wherein the bandwidth request in the receiving step is a volume request and is received over at least one of a contention channel and a data channel, the computer-readable medium further comprising computer-executable instructions for causing the computer system to perform the steps of:

receiving a piggybacked volume request over the data channel;

placing the piggybacked volume request in a corresponding one of the global queues;

determining whether the plurality of channels are oversubscribed; and selectively discarding the piggybacked volume request based upon the step of determining whether the plurality of channels are oversubscribed.

Mon

43. (Original) The computer-readable medium as in claim 42, wherein the step of determining whether the plurality of channels are oversubscribed comprises:

determining whether each of the plurality of local queues exceeds a respective queuing threshold.

44. (Original) The computer-readable medium as in claim 35, wherein the plurality of channels are designated as data channels that are sequentially ordered, the allocating step comprising:

selectively assigning the transmission slots according to a prescribed order of the data channels based upon the bandwidth request type.

- 45. (Original) The computer-readable medium as in claim 44, wherein the prescribed order in the selectively assigning step begins with the first data channel if the bandwidth request type is rate request.
- 46. (Original) The computer-readable medium as in claim 44, wherein the prescribed order in the selectively assigning step begins with the last data channel if the bandwidth request type is volume request.
- 47. (Original) The computer-readable medium as in claim 35, further comprising computer-executable instructions for causing the computer system to perform the steps of:

receiving a plurality of rate requests;

A M

receiving a defragmentation command; and

moving the rate requests from the local queues to the corresponding global queues for reallocation in response to the defragmentation command.

48. (Original) The computer-readable medium as in claim 35, wherein the bandwidth request is an original volume request, the computer-readable medium further comprising computer-executable instructions for causing the computer system to perform the steps of:

receiving a follow-up volume request;

associating the follow-up volume request with the original volume request; and placing the follow-up volume request to a particular local queue that stored the original bandwidth request among the plurality of local queues.

49. (Original) The computer-readable medium as in claim 48, wherein the associating step comprises:

maintaining a database of pointers for the terminal, one of the pointers specifying the particular local queue.

50. (Original) The computer-readable medium as in claim 35, further comprising computer-executable instructions for causing the computer system to perform the steps of:

receiving a plurality of volume requests; and spreading the volume requests across each of the local queues.

A A

Alak

51. (Original) The computer-readable medium as in claim 50, wherein each of the local queues has a counter that counts a quantity of the volume requests in the respective local queue, the distributing step comprising:

comparing counter values of the counters with respective predetermined thresholds corresponding to the local queues.